Lecture 4

Structured Query Language (basic)

COMP3278A

Introduction to Database Management Systems

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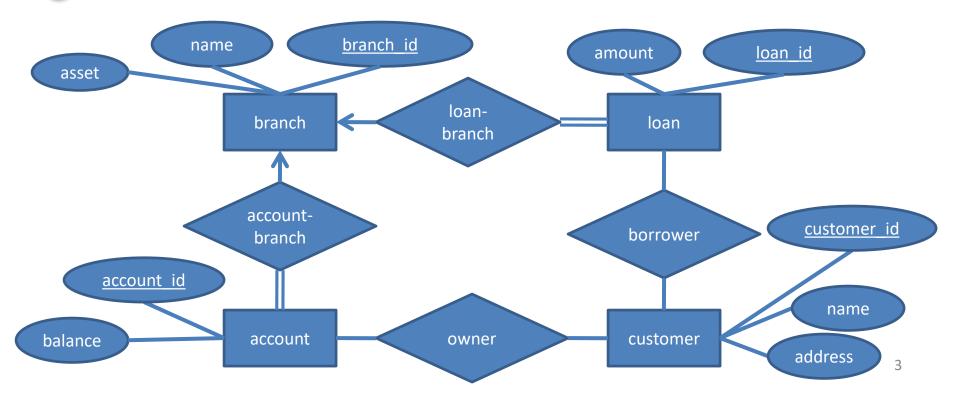
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Outcome based learning (OBL)

- Outcome 1. Information Modeling
 - Able to understand the modeling of real life information in a database system.
- Outcome 2. Query Languages
 - Able to understand and use the languages designed for data access.
- Outcome 3. System Design
 - Able to understand the design of an efficient and reliable database system.
- Outcome 4. Application Development
 - Able to implement a practical application on a real database.

Recap

- Let's consider the following steps in developing a database application in a banking enterprise.
- Step 1. Information modeling



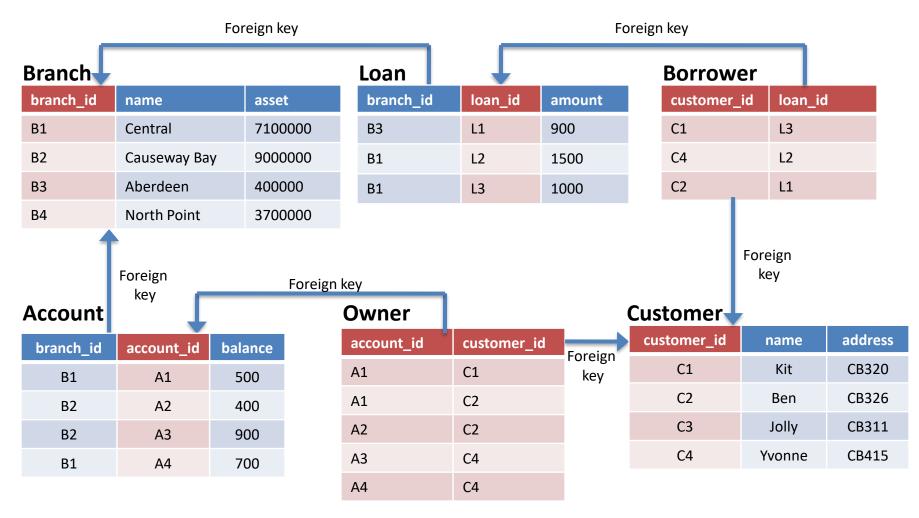
Recap

Step 2. Reduce to database table definitions

Recap

- Step 3. Create the database and tables
- Step 4. Design the SQL to access data for the application
- Step 5. Relational Algebra optimizes SQL (DBMS does this automatically). We'll learn its basic concepts.

Running example



What is SQL?

- Structured Query Language (pronounced as "sequel")
- Language for defining, modifying and querying data in an RDBMS.
- SQL is declarative
 - Concerns about the task we want to accomplish, without specifying how.
- SQL has many standards and implementations
 - Read the documentation on which features are supported exactly.

Section 1

Create and Drop Table

Create table

A database table is defined using the CREATE TABLE command.

Table name (cannot be a keyword in database, e.g., **Create**)

Column name

Constraint (Primary Key), optional in this statement, can add back later using the ALTER TABLE command.

PRIMARY KEY
automatically ensures
NOT NULL.

create table Branch
(
branch_id VARCHAR(15) NOT
NULL,
name VARCHAR(30) NOT NULL,
asset INT UNSIGNED NOT NULL,
PRIMARY KEY(branch_id)
);

No comma in the last instruction.

Column type

NOT NULL means each record's value in the column must not be a null value.

Drop table

DROP TABLE deletes all information about the dropped table from the database.

DROP TABLE Branch;

The DBMS may reject the DROP TABLE instruction when the table is referenced by another table via some constraints (e.g., referential constraints).

(Customer	•		Borrower		
	customer_id	name	address	customer_id	loan_id	
	C1	Kit	CB320	C1	L3	
	C2	Ben	CB326	C4	L2	
	C3	Jolly	CB311	C2	L1	
	CA	Vyonno	CD/11E			

After the foreign key is established, if we drop the Customer table, the records in the Borrower table will lost their references.

(i.e., Cannot find out who borrow the loan anymore.)



Alter table

- ALTER TABLE can be used to
 - Add columns to an existing table.

ALTER TABLE Branch **ADD** branch_phone **INT (12)**;

Remove a column from a table.

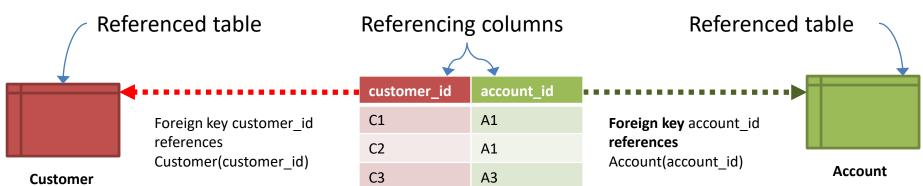
ALTER TABLE Branch **DROP** branch_phone;

Add constraints (e.g., PRIMARY KEY) to a table.

ALTER TABLE Branch **ADD PRIMARY KEY** (branch_id);

Foreign key constraints

- A foreign key is a referential constraint between two tables.
- The columns in the referencing table must reference the columns of the primary key or other superkey in the referenced table.
 - i.e., The value in one row of the referencing columns must occur in a single row in the referenced table. The referencing columns must be primary/candidate key of another table. The referencing table cannot contain record that doesn't exist in the referenced table.



Owner

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Foreign key constraints

The foreign key can be established in the CREATE TABLE command.

```
CREATE TABLE Owner
(
    customer_id VARCHAR(15),
    account_id VARCHAR(15),
    PRIMARY KEY(customer_id, account_id),
    FOREIGN KEY(customer_id) REFERENCES Customer(customer_id),
    FOREIGN KEY(account_id) REFERENCES Account(account_id)
);
```

The foreign key can also be defined using the ALTER TABLE command.

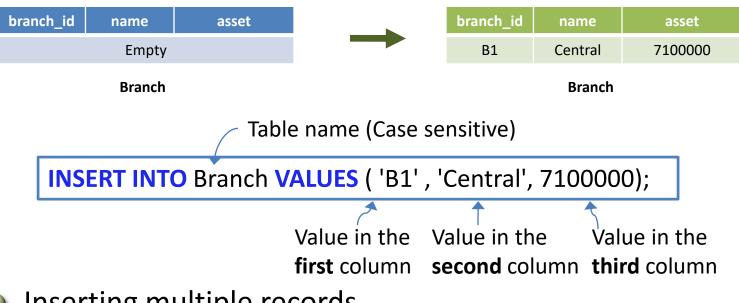
```
ALTER TABLE Owner
ADD FOREIGN KEY (customer_id) REFERENCES Customer(customer_id);
```

Section 2

Insert, Delete and Update

The INSERT clause

The INSERT INTO command is used to insert records (tuples) into the database table.



Inserting multiple records

```
INSERT INTO Branch VALUES
('B2', 'Causeway Bay', 9000000),
('B3', 'Aberdeen', 400000);
```

The INSERT clause

Most DBMS provide an alternative way to insert large amount of records into a table.

E.g., LOAD DATA LOCAL INFILE in MySQL. MySQL.

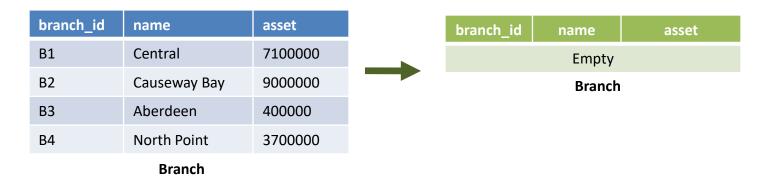
LOAD DATA LOCAL INFILE 'text.txt'
INTO TABLE Branch
FIELDS TERMINATED BY ';'
LINES TERMINATED BY '\n';

B1;Central;7100000 B2;Causeway Bay;9000000 B3;Aberdeen;400000 B4; North Point;3700000

text.txt

The DELETE clause

- The DELETE FROM command is used to delete records (tuples) from a database table.
 - Query: Delete all records from the Branch table.



DELETE FROM Branch;

The DELETE clause

- The DELETE FROM command is used to delete records (tuples) from a database table.
 - Query: Delete the branch "Central" from the Branch table.

branch_id	name	asset
B1	Central	7100000
	Causeway Bay	9000000
3	Aberdeen	400000
B4	North Point	3700000
	Branch	

DELETE FROM Branch WHERE name = 'Central';

The tuples that satisfy the conditions specified here are deleted.

- The UPDATE command is used to update records (tuples) from a database table.
 - Query: Update the asset of branch with branch_id 'B1' to \$0.

branch_id	name	asset		branch_id	name
B1	Central	7100000		B1	Central
B2	Causeway Bay	9000000	\rightarrow	B2	Causeway Bay
В3	Aberdeen	400000		В3	Aberdeen
B4	North Point	3700000		B4	North Point
Branch					Branch

UPDATE Branch
SET asset = 0
WHERE branch_id = 'B1';

- The UPDATE command can also be used with arithmetic expressions.
 - Query: Increase all accounts with balances over \$500 by 6%.

account_id	branch_id	balance
A1	B1	500
A2	B2	400
A3	B2	900
A4	B1	700
	Account	

UPDATE Account **SET** balance = balance * 1.06 **WHERE** balance > 500;

- The UPDATE command can also be used with arithmetic expressions.
 - Query: Increase all accounts with balances under \$500 by 5% and all other accounts by 6%.

account_id	branch_id	balance		account_id	branch_id
A1	B1	500		A1	B1
A2	B2	400	—	A2	B2
A3	B2	900		A3	B2
A4	B1	700		A4	B1

Account Account

SET balance = balance * 1.05 **WHERE** balance < 500; UPDATE Account
SET balance = balance * 1.06
WHERE balance >= 500;



The CASE command can be used to perform conditional update.

```
UPDATE Account

SET balance = CASE

WHEN balance <=500 THEN balance *1.05

ELSE balance * 1.06

END
```

Note: When there are multiple **WHEN ... THEN** in the query, only the first true statement (from top to bottom) will be executed.

Section 3

Querying

The SELECT clause

- The SELECT clause lists the attributes desired in the result of a query.
 - Query: Find the names of all customers.

customer_id	name	address		name
C1	Kit	CB320		Kit
C2	Ben	CB326	\rightarrow	Ben
C3	Jolly	CB311		Jolly
C4	Yvonne	CB415		Yvonne
	Customer			Result

SELECT name **FROM** Customer;

- An asterisk in the select clause denotes "all attributes"
 - Query: List all column values of all customer records.

SELECT * **FROM** Customer;

The SELECT clause

- The SELECT clause can contain arithmetic expressions (+, -, *, /) operating on constants or attributes of tuples.
 - Query: List the loan_id and amount of each loan record, display the amount in USD (originally stored in HKD).

loan_id	branch_id	amount
L1	В3	900
L2	B2	1500
L3	B1	1000
	Loan	

SELECT loan_id, amount/7.8 **FROM** Loan;

- The FROM clause lists the relations (tables) involved in the query.
 - Query: Find the Cartesian product of Customer and Borrower

SELECT * **FROM** Customer, Borrower;

customer_id	name	address	customer_id	loan_id

customer_id	name	address
C1	Kit	CB320
C2	Ben	CB326
С3	Jolly	CB311
C4	Yvonne	CB415

Customer

customer_id	loan_id
C1	L3
C4	L2
C2	L1





Cartesian product of A and B means generate **all possible pairs** of records from A and B.

- The FROM clause lists the relations (tables) involved in the query.
 - Query: Find the Cartesian product of Customer and Borrower

SELECT * **FROM** Customer, Borrower;

customer_id	name	address	customer_id	loan_id
C1	Kit	CB320	C1	L3

customer_id	name	address			
C1	Kit	CB320			
C2	Ben	CB326			
C3	Jolly	CB311			
C4	Yvonne	CB415			
Customer					

customer_id	loan_id
C1	L3
C4	L2
C2	L1



Cartesian product of A and B means generate all possible pairs of records from A and B.

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SELECT * **FROM** Customer, Borrower;

customer_id	name	address	customer_id	loan_id
C1	Kit	CB320	C1	L3
C2	Ben	CB326	C1	L3

customer_id	name	address	
C1	Kit	CB320	
C2	Ben	CB326	
C3	Jolly	CB311	
C4	Yvonne	CB415	
Customer			

customer_id	loan_id
C1	L3
C4	L2
C2	L1



Cartesian product of A and B means generate all possible pairs of records from A and B.

- The FROM clause lists the relations (tables) involved in the query.
 - Query: Find the Cartesian product of Customer and Borrower

SELECT * **FROM** Customer, Borrower;

Cartesian product is the most primitive way of joining two tables. However, many resulting tuples are not very useful. Therefore, we often need to specify the joining condition to filter out the non-meaningful results.



customer_id	name	address	customer_id	loan_id
C1	Kit	CB320	C1	L3
C2	Ben	CB326	C1	L3
C3	Jolly	CB311	C1	L3
C4	Yvonne	CB415	C1	L3
C1	Kit	CB320	C4	L2
C2	Ben	CB326	C4	L2
C3	Jolly	CB311	C4	L2
C4	Yvonne	CB415	C4	L2
C1	Kit	CB320	C2	L1
C2	Ben	CB326	C2	L1
C3	Jolly	CB311	C2	L1
C4	Yvonne	CB415	C2	L1

- The WHERE clause specifies conditions that the result must satisfy.
 - Query: For each loan, find out the name of the customer who borrow the loan.

Let us learn the process of constructing the SQL for this query.



The WHERE clause specifies conditions that the result must satisfy.

Query: For each loan, find out the name of the customer who borrow the loan.

customer_id	loan_id
C1	L3
C4	L2
C2	L1

Borrower

customer_id	name	address
C1	Kit	CB320
C2	Ben	CB326
C3	Jolly	CB311
C4	Yvonne	CB415

Customer

Step 1. What are the table(s) that contain the information to answer this query?



Observation 1.

First, the information of customers (customer_id) who borrow loan is in the **Borrower** table.



Observation 2.

Second, we need to find out the name of the customer, the name is in the **Customer** table.



SELECT Borrower.loan_id, Customer.name FROM Customer, Borrower

Step 2. Now we want to relate two tables, if no conditions is specified, Cartesian product will be returned. What is the joining condition?



customer_id	loan_id
C1	L3
C4	L2
C2	L1

Borrower



customer_id	name	address
C1	Kit	CB320
C2	Ben	CB326
C3	Jolly	CB311
C4	Yvonne	CB415

Customer

customer_id	name	address	customer_id	loan_id
C1	Kit	CB320	C1	L3
C2	Ben	CB326	C1	L3
C3	Jolly	CB311	C1	L3
C4	Yvonne	CB415	C1	L3
C1	Kit	CB320	C4	L2
C2	Ben	CB326	C4	L2
C3	Jolly	CB311	C4	L2
C4	Yvonne	CB415	C4	L2
C1	Kit	CB320	C2	L1
C2	Ben	CB326	C2	L1
C3	Jolly	CB311	C2	L1
C4	Yvonne	CB415	C2	L1

SELECT Borrower.loan_id, Customer.name

FROM Customer, Borrower

WHERE Customer.customer_id =

Borrower.customer_id

loan_id	name
L3	Kit
L2	Yvonne
L1	Ben





customer_id	loan_id
C1	L3
C4	L2
C2	L1

Borrower



customer_id	name	address	
C1	Kit	CB320	
C2	Ben	CB326	
C3	Jolly	CB311	
C4	Yvonne	CB415	

Customer

customer_id	name	address	customer_id	loan_id	
C1	Kit	CB320	C1	L3	
C2	Ben	CB326	C1	L3	
C3	Jolly	CB311	C1	L3	
C4	Yvonne	CB415	C1	L3	
C1	Kit	CB320	C4	L2	
C2	Ben	CB326	C4	L2	
C3	Jolly	CB311	C4	L2	
C4	Yvonne	CB415	C4	L2	
C1	Kit	CB320	C2	L1	
C2	Ben	CB326	C2	L1	
C3	Jolly	CB311	C2	L1	
C4	Yvonne	CB415	C2	L1	
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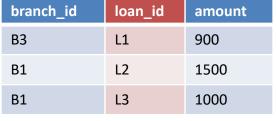
The WHERE clause specifies conditions that the result must satisfy.

Comparison results can be combined using logical connectives AND, OR, and NOT.

Query: Find all loan ID of loans made at branch_id B1 with

loan amounts >\$1200.

	There are two
	conditions in the query!





Loan

```
SELECT loan_id

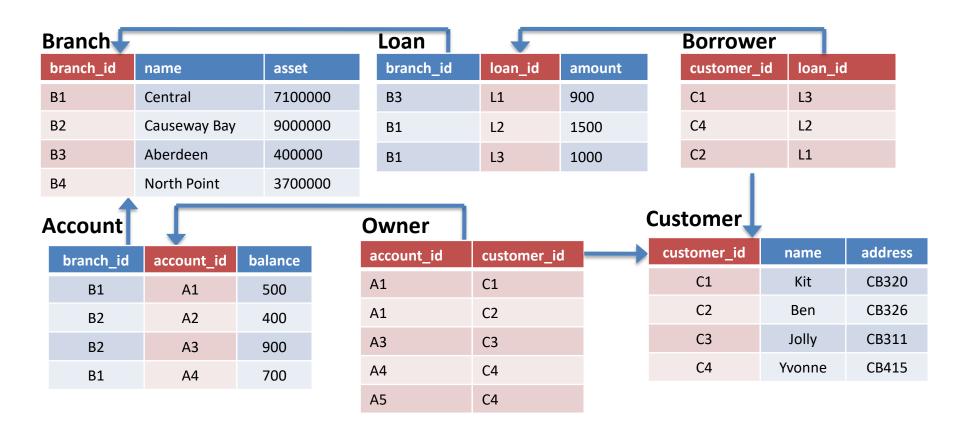
FROM Loan

WHERE branch_id = 'B1' AND

amount > 1200;
```

Exercise

Query: Find the names of all branches that have a loan.



Exercise

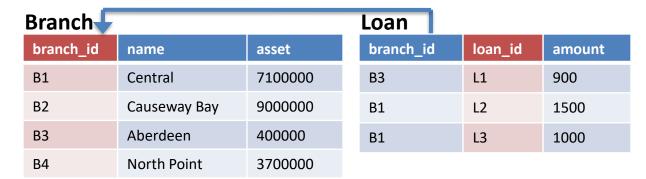
- Query: Find the names of all branches that have a loan.
 - Step 1. Identify the tables that contain the necessary information to answer the query.

Branch		Loan	oan		
branch_id	name	asset	branch_id	loan_id	amount
B1	Central	7100000	В3	L1	900
B2	Causeway Bay	9000000	B1	L2	1500
В3	Aberdeen	400000	B1	L3	1000
B4	North Point	3700000			

Step 2. Construct the SELECT statement.

```
SELECT ?
FROM Branch, Loan
WHERE ?
;
```

- Query: Find the names of all branches that have a loan.
 - Step 1. Identify the tables that contain the necessary information to answer the query.



Step 2. Construct the SELECT statement.

```
SELECT Branch.name
FROM Branch, Loan
WHERE?
;
```

Query: Find the names of all branches that have a loan.

Step 1. Identify the tables that contain the necessary information to answer the query.

Branch			Loan	_	
branch_id	name	asset	branch_id	loan_id	amount
B1	Central	7100000	В3	L1	900
B2	Causeway Bay	9000000	B1	L2	1500
В3	Aberdeen	400000	B1	L3	1000
B4	North Point	3700000			

Step 2. Construct the SELECT statement.

```
SELECT Branch.name
FROM Branch. Loan
WHERE Branch.branch_id = Loan.branch_id
;
```

Usually, when linking the information of two tables, we need to specify the joining condition. Often we need to join the columns that participate in the referential constraint between the two tables.

Joining condition

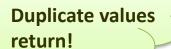


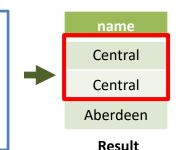
- Query: Find the names of all branches that have a loan.
 - Step 1. Identify the tables that contain the necessary information to answer the query.

Branch			Loan		
branch_id	name	asset	branch_id	loan_id	amount
B1	Central	7100000	В3	L1	900
B2	Causeway Bay	9000000	B1	L2	1500
В3	Aberdeen	400000	B1	L3	1000
R/I	North Point	370000			

Step 2. Construct the SELECT statement.

SELECT Branch.name
FROM Branch, Loan
WHERE Branch.branch_id = Loan.branch_id
;





Query: Find the names of all branches that have a loan.

Step 1. Identify the tables that contain the necessary

information to answer the query.

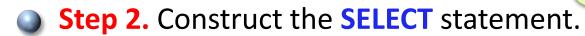
Branch	Loan

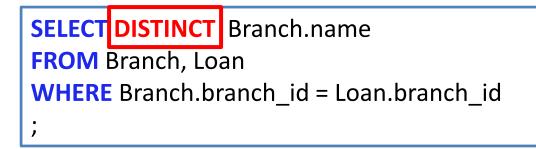
branch_id	name	asset
B1	Central	7100000
B2	Causeway Bay	9000000
В3	Aberdeen	400000
B4	North Point	3700000

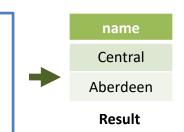
branch_id	loan_id	amount
В3	L1	900
B1	L2	1500
B1	L3	1000

You can eliminate duplicate values in the results by using the **DISTINCT** keyword.

Duplicate values return!







Section 4

Renaming

Renaming

```
SELECT DISTINCT Branch.name

FROM Branch, Loan

WHERE Branch.branch_id = Loan.branch_id

;
```

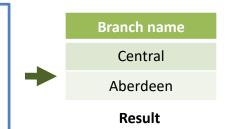
- Rename can be operated on both tables and attributes.
 - Rename on attribute.

I want to **rename** the column "name" in the result into "Branch name".



We use the keyword AS to signify renaming.

SELECT DISTINCT Branch.name AS 'Branch name'
FROM Branch, Loan
WHERE Branch.branch_id = Loan.branch_id
;



name

Central

Aberdeen

Renaming

```
SELECT DISTINCT Branch.name

FROM Branch, Loan

WHERE Branch.branch_id = Loan.branch_id

Result
```

- Rename can be operated on both tables and attributes.
 - Rename on tables.

The two SQLs are equivalent to each other.



```
SELECT DISTINCT B.name

FROM Branch B, Loan L

WHERE B.branch_id = L.branch_id

Result
```

Section 5

String operations

The LIKE clause

- The most commonly used operation on strings is pattern matching using LIKE.
 - Percent(%):matches any substring.
 - Underscore(_): matches any character.
 - 'Perry%' matches any string beginning with "Perry".
 - \(\scale= \sc
- Note: Patterns are case sensitive.

The LKE clause

Query: Find the names of all customers whose address includes the substring '320'.

Customer

customer_id	name	address
C1	Kit	CB320
C2	Ben	CB326
C3	Jolly	CB311
C4	Yvonne	CB415

SELECT name **FROM** Customer **WHERE** address **LIKE** '%320%';

https://dev.mysql.co m/doc/refman/8.0/e n/patternmatching.html

name

Kit

Result

Question: How about matching using regular expression? ©



The LIKE clause

WHERE name LIKE 'a%'	Finds any values that start with "a"
WHERE name LIKE '%a'	Finds any values that end with "a"
WHERE name LIKE '%or%'	Finds any values that have "or" in any position
WHERE name LIKE '_r%'	Finds any values that have "r" in the second position
WHERE name LIKE 'a%'	Finds any values that start with "a" and are at least 3 characters in length
WHERE name LIKE 'a%o'	Finds any values that start with "a" and ends with "o"

Section 6

Ordering results

The ORDER BY clause

- The ORDER BY clause list the result in sorted order.
 - Query: List the names of all customers in alphabetic order.

Customer

customer_id	name	address		name
C1	Kit	CB320		Ben
C2	Ben	CB326	\rightarrow	Jolly
C3	Jolly	CB311		Kit
C4	Yvonne	CB415		Yvonne
				Result

SELECT name
FROM Customer
ORDER BY name ASC;

Yvonne

Kit

Jolly

Ben

SELECT name **FROM** Customer **ORDER BY** name **DESC**;

The ORDER BY clause

- The ORDER BY clause list the result in sorted order.
 - Query: List the loan records in ascending order of the branch_id, if two tuples having the same branch_id, order by their loan amount in descending order.



Intermediate Result

Final Result

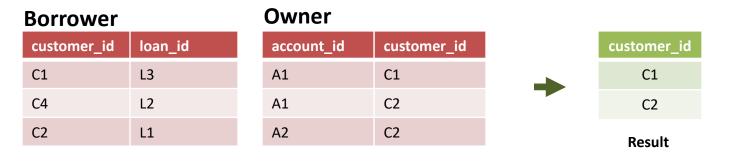
SELECT *
FROM Loan
ORDER BY branch_id ASC,
amount DESC;

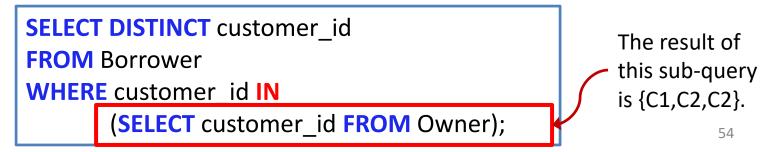
Section 7

Simple Nested Query

The N clause

- The IN clause allows you to specify discrete values in the WHERE search criteria.
 - Query: Find the customer id of all customers who have both an account and a loan.



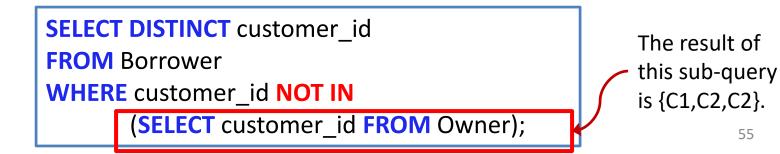


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The N clause

- The IN clause allows you to specify discrete values in the WHERE search criteria.
 - Query: Find the customer_id of all customers who have a lone but not having an account.

Owner Borrower customer id customer id loan id account id customer_id A1 C1 L3 C4 C4 L2 Α1 C2 C2 L1 A2 C2 Result



Section 8

Aggregation

Aggregate functions

- Aggregation functions take a collection of values as input and return a single value.
 - Query: Find the average balance of all accounts at the branch with branch_id 'B2'.

Account

branch_id	account_id	balance
B1	A1	500
В2	A2	400
B2	A3	900
B1	A4	700

SELECT AVG(balance)
FROM Account
WHERE branch_id = 'B2';

Aggregate functions

- Aggregation functions.
 - AVG
 - MIN
 - MAX
 - SUM
 - COUNT

The GROUP BY clause

- Aggregation function can be applied to a group of sets of tuples by using GROUP BY clause.
 - Query: Find the average balance at each branch.

Step1. Grouping GROUP BY branch_id

Account

branch_id	account_id	balance		branch_id	account_id	balance
B1	A1	500	—	B1	A1	500
B2	A2	400			A4	700
B2	A3	900		D2	A2	400
B1	A4	700		B2	A3	900

The GROUP BY clause

- Aggregation function can be applied to a group of sets of tuples by using GROUP BY clause.
 - Query: Find the average balance at each branch.

SELECT branch_id, **AVG**(balance) **FROM** Account **GROUP BY** branch_id;

Step1. Grouping **GROUP BY branch_id**

Step2. Aggregation AVG(balance)

Account

branch_id	account_id	balance
B1	A1	500
B2	A2	400
B2	A3	900
B1	A4	700

branch_id	account_id	balance	
B1	A1	500	
	A4	700	-
D2	A2	400	
B2	А3	900	

branch_id	AVG (balance)
B1	600.0000
B2	650.0000
Result	

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The HAVING clause

- It is useful to state a condition that applies to groups rather than to tuples.
 - Query: Find the branches where the average account balance is no less than \$650.

SELECT branch_id, AVG(balance)
FROM Account
GROUP BY branch_id
HAVING AVG(balance) >= 650;

	branch_id	AVG (balance)
	B2	650.0000
	R	esult
Step3. Filtering Having AVG(balance) >= 650)		

Account

branch_id	account_id	balance	
B1	A1	500	
B2	A2	400	\rightarrow
B2	A3	900	
B1	A4	700	

branch_id	account_id	balance	
D1	A1	500	
B1	A4	700	٠
D2	A2	400	
B2	A3	900	

branch_id	AVG (balance)
B1	600.0000
B2	650.0000

Section 9

Join

Join

A join takes 2 tables as input and returns a table.

Employee

e_name	department_id
c_name	acpartment_ia
Kit	31
Ben	33
John	33
Jolly	34
Yvonne	34
David	NULL

Department

department_id	d_name
31	CS
33	Civil
34	ME
35	EEE

Cartesian product, then
E.department_id = D.department_id

e_name	E.department_id	D.department_id	d_name
Kit	31	31	CS
Ben	33	33	Civil
John	33	33	Civil
Jolly	34	34	ME
Yvonne	34	34	ME

The OUTER JOIN clause

An outer join does not require each record in the two joined tables to have a matching record.

Employee

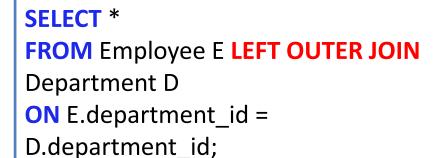
e_name	department_id
Kit	31
Ben	33
John	33
Jolly	34
Yvonne	34
David	NULL

Department

department_id	d_name
31	CS
33	Civil
34	ME
35	EEE



Even if the LEFT table record does not have matching records in the RIGHT table, we still output the tuple in the LEFT table (with null values for the columns of the RIGHT table).



e_name	E.department_id	D.department_id	d_name
Kit	31	31	CS
Ben	33	33	Civil
John	33	33	Civil
Jolly	34	34	ME
Yvonne	34	34	ME
David	NULL	NULL	NULL

The OUTER JOIN clause

An outer join does not require each record in the two joined tables to have a matching record.

Employee

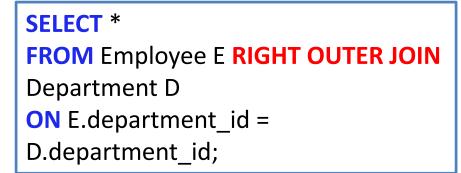
e_name	department_id
Kit	31
Ben	33
John	33
Jolly	34
Yvonne	34
David	NULL

Department

department_id	d_name
31	CS
33	Civil
34	ME
35	EEE



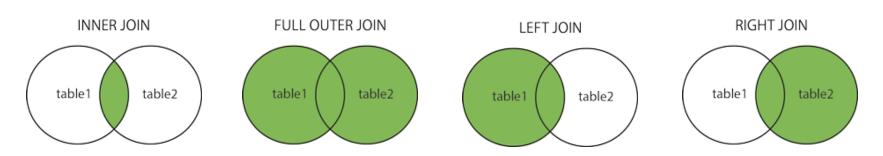
Even if the RIGHT table record does not have matching records in the LEFT table, we still output the tuple in the RIGHT table (with null values for the columns of the LEFT table).



e_name	E.department_id	D.department_id	d_name
Kit	31	31	CS
Ben	33	33	Civil
John	33	33	Civil
Jolly	34	34	ME
Yvonne	34	34	ME
NULL	NULL	35	EEE

Here are the different types of the JOINs in SQL:

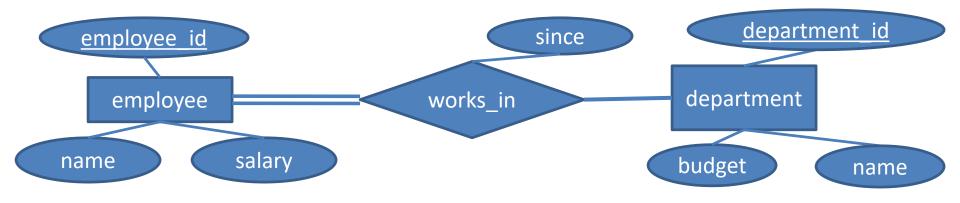
- •(INNER) JOIN: Returns records that have matching values in both tables
- •LEFT (OUTER) JOIN: Returns all records from the left table, and the matched records from the right table
- •RIGHT (OUTER) JOIN: Returns all records from the right table, and the matched records from the left table
- •FULL (OUTER) JOIN: Returns all records when there is a match in either left or right table



Section 10 More on SQL

- An Example
- Set operations
- More on nested queries
- Null values
- Views
- Authorization
- Assertion
- Other SQL constructs

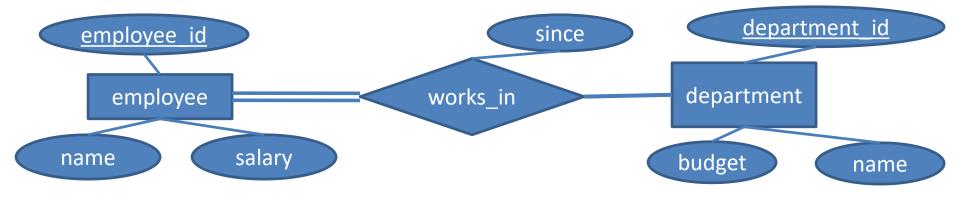
Step 1. Information modeling.



Before we proceed to learning more SQL constructs, lets have a revision on what we have learned up to this chapter.



Step 1. Information modeling



- Step 2. Reduce to database tables
 - Employees (employee id, name, salary)
 Foreign key: none.
 - Departments (<u>department id</u>, name, budget) Foreign key : none.
 - Works_in(employee id, department id, since)
 Foreign key: employee_id REFERENCES Employee (employee_id).
 department_id REFERENCES Department (department_id).

Step 3. Create the database

```
CREATE TABLE Employees (
  employee id INT(12),
                                            INNODB storage engine,
  name VARCHAR(30) NOT NULL,
                                            just for MySQL to support
  salary INT UNSIGNED NOT NULL,
                                            foreign key constraints.
  PRIMARY KEY(employee id)
)ENGINE = INNODB: ◀
CREATE TABLE Departments (
  department id INT(12),
  name VARCHAR(30) NOT NULL,
  budget INT UNSIGNED NOT NULL,
  PRIMARY KEY(department id)
) ENGINE = INNODB;
CREATE TABLE Works in(
  employee id INT(12),
  department id INT(12),
  since DATE NOT NULL,
  PRIMARY KEY(employee_id, department_id),
  FOREIGN KEY (employee id) REFERENCES Employees (employee id),
  FOREIGN KEY (department id) REFERENCES Departments (department id)
) ENGINE = INNODB;
```



Step 3. Create the database

```
INSERT INTO Employees VALUES (1, 'Jones', 26000);
INSERT INTO Employees VALUES (2, 'Smith', 28000);
INSERT INTO Employees VALUES (3, 'Parker', 35000);
INSERT INTO Employees VALUES (4, 'Smith', 24000);
```

INSERT INTO Departments VALUES (1, 'Toys', 122000), (2, 'Tools', 239000), (3, 'Food', 100000);

INSERT INTO Works in **VALUES** (1, 1, '2001-1-1'), (2, 1, '2002-4-1'), (2, 2, '2005-2-2'), (3, 3, '2003-1-1'), (4, 3, '2005-1-1');

Works in

3

3

since

2003-1-1

2005-1-1

Employees

name

Jones

Smith

Parker

Smith

salary

26000

28000

35000

24000

employee_id

1

2

3

4

employee id department_id 1 2001-1-1 1 2002-4-1 2 1 2 2 2005-2-2

3

4

department_id	name	budget
1	Toys	122000
2	Tools	239000
3	Food	100000

- Step 4. Design the SQL to access data for the application
 - Query 1: Find the names of all employees and remove duplicates.

Employees

name	salary
Jones	26000
Smith	28000
Parker	35000
Smith	24000
	Jones Smith Parker

Works in

employee_id	department_id	since
1	1	2001-1-1
2	1	2002-4-1
2	2	2005-2-2
3	3	2003-1-1
4	3	2005-1-1

department_id	name	budget
1	Toys	122000
2	Tools	239000
3	Food	100000

Query 2: Find the employee_ids and names of employees who work in department with department_id=2.

Employees

employee_id	name	salary
1	Jones	26000
2	Smith	28000
3	Parker	35000
4	Smith	24000

Works in

employee_id	department_id	since
1	1	2001-1-1
2	1	2002-4-1
2	2	2005-2-2
3	3	2003-1-1
4	3	2005-1-1

name	budget
Toys	122000
Tools	239000
Food	100000
	Toys Tools

Query 3: Find the dept. names where employee with employee_id = 2 works.

Employees

employee_id	name	salary
1	Jones	26000
2	Smith	28000
3	Parker	35000
4	Smith	24000

Works_in

employee_id	department_id	since
1	1	2001-1-1
2	1	2002-4-1
2	2	2005-2-2
3	3	2003-1-1
4	3	2005-1-1

department_id	name	budget
1	Toys	122000
2	Tools	239000
3	Food	100000

Query 4: Find the dept. ids where employees named Smith work.

Employees

employee_id	name	salary
1	Jones	26000
2	Smith	28000
3	Parker	35000
4	Smith	24000

Works_in

employee_id	department_id	since
1	1	2001-1-1
2	1	2002-4-1
2	2	2005-2-2
3	3	2003-1-1
4	3	2005-1-1

department_id	name	budget
1	Toys	122000
2	Tools	239000
3	Food	100000

Query 5: Find the dept. names where employees named Smith work.

Employees

employee_id	name	salary
1	Jones	26000
2	Smith	28000
3	Parker	35000
4	Smith	24000

Works_in

employee_id	department_id	since
1	1	2001-1-1
2	1	2002-4-1
2	2	2005-2-2
3	3	2003-1-1
4	3	2005-1-1

department_id	name	budget
1	Toys	122000
2	Tools	239000
3	Food	100000

Query 6: Find the names of departments which have an employee named Smith and their budget is greater than 100000.

Employees

employee_id	name	salary
1	Jones	26000
2	Smith	28000
3	Parker	35000
4	Smith	24000

Works_in

employee_id	department_id	since
1	1	2001-1-1
2	1	2002-4-1
2	2	2005-2-2
3	3	2003-1-1
4	3	2005-1-1

department_id	name	budget
1	Toys	122000
2	Tools	239000
3	Food	100000

Query 7: Find the budgets of departments, who employ an employee called 'Smith'.

Employees

employee_id	name	salary
1	Jones	26000
2	Smith	28000
3	Parker	35000
4	Smith	24000

Works_in

employee_id	department_id	since
1	1	2001-1-1
2	1	2002-4-1
2	2	2005-2-2
3	3	2003-1-1
4	3	2005-1-1

department_id	name	budget
1	Toys	122000
2	Tools	239000
3	Food	100000
2		

Query 8: For each department, find the total number of employees it employs.

Employees

name	salary
Jones	26000
Smith	28000
Parker	35000
Smith	24000
	Jones Smith Parker

Works_in

employee_id	department_id	since
1	1	2001-1-1
2	1	2002-4-1
2	2	2005-2-2
3	3	2003-1-1
4	3	2005-1-1

department_id	name	budget
1	Toys	122000
2	Tools	239000
3	Food	100000

Query 9: Find the dept. names with at least 2 employee.

Employees

employee_id	name	salary
1	Jones	26000
2	Smith	28000
3	Parker	35000
4	Smith	24000

Works_in

employee_id	department_id	since
1	1	2001-1-1
2	1	2002-4-1
2	2	2005-2-2
3	3	2003-1-1
4	3	2005-1-1

department_id	name	budget
1	Toys	122000
2	Tools	239000
3	Food	100000

Query 10: Find the employee_id of all employees whose name includes the substring "one".

Employees

employee_id	name	salary
1	Jones	26000
2	Smith	28000
3	Parker	35000
4	Smith	24000

Works_in

employee_id	department_id	since
1	1	2001-1-1
2	1	2002-4-1
2	2	2005-2-2
3	3	2003-1-1
4	3	2005-1-1

Departments

department_id	name	budget
1	Toys	122000
2	Tools	239000
3	Food	100000

Query 11: Find the employee_id and name of the employees who worked in the departments with budget more than 100,000.

Employees

employee_id	name	salary
1	Jones	26000
2	Smith	28000
3	Parker	35000
4	Smith	24000

Works_in

employee_id	department_id	since
1	1	2001-1-1
2	1	2002-4-1
2	2	2005-2-2
3	3	2003-1-1
4	3	2005-1-1

Departments

department_id	name	budget
1	Toys	122000
2	Tools	239000
3	Food	100000

Query 12: Find the name and budget of the department with the greatest budget.

Employees

employee_id	name	salary
1	Jones	26000
2	Smith	28000
3	Parker	35000
4	Smith	24000

Works_in

employee_id	department_id	since
1	1	2001-1-1
2	1	2002-4-1
2	2	2005-2-2
3	3	2003-1-1
4	3	2005-1-1

Departments

department_id	name	budget
1	Toys	122000
2	Tools	239000
3	Food	100000

Query 13: Find the names of employees who work in at least 2 departments.

Employees

employee_id	name	salary
1	Jones	26000
2	Smith	28000
3	Parker	35000
4	Smith	24000

Works_in

employee_id	department_id	since
1	1	2001-1-1
2	1	2002-4-1
2	2	2005-2-2
3	3	2003-1-1
4	3	2005-1-1

Departments

department_id	name	budget
1	Toys	122000
2	Tools	239000
3	Food	100000

Query 14: In each department, find the highest salary of the employee in that department.

Employees

employee_id	name	salary
1	Jones	26000
2	Smith	28000
3	Parker	35000
4	Smith	24000

Works_in

employee_id	department_id	since
1	1	2001-1-1
2	1	2002-4-1
2	2	2005-2-2
3	3	2003-1-1
4	3	2005-1-1

department_id	name	budget
1	Toys	122000
2	Tools	239000
3	Food	100000

Q&A Game

Please read the below table schemas and answer four SQL questions.

- Restaurant (<u>restaurant id</u>, name)
- RestaurantCategory (<u>restaurant_id</u>, <u>category</u>)
 - Foreign key: {restaurant_id} references Restaurant
- **Branch** (*restaurant id*, *branch no*, *location*, *seats*)
 - Foreign key: {restaurant_id} references Restaurant
- Member (<u>member id</u>, name, birthday, joined, points)
 - (joined is the year the member joined)
- Visits (visit_id, member_id, restaurant_id, branch_no, date, score)
 - Foreign keys:{restaurant_id, branch_no} references Branch, {member_id} references Member

Restaurant

restaurant_id	name
1	McRonalds
2	PizzaHub
3	DeliItaly
4	Ultra Sandwich
5	Starducks

RestaurantCategory

restaurant_	_id category
1	Fast Food
1	Take Away
2	Italian
3	Cafe
3	Italian
4	Light meal
5	Cafe
5	Light meal
5	Take Away

Member

member_	id name	birthday	joined	points
1	Cleo	1983-09-25	2017	690
2	Evan	1988-03-28	1989	130
3	Todd	1966-06-22	1967	190
4	Lonny	1973-04-05	2016	10
5	Keith	1991-06-19	1992	380
6	Royce	1965-06-14	2006	300
7	Mavis	1977-08-21	1981	840
8	Alvin	1998-04-17	2008	900
9	Ira	1993-07-17	2015	100
10	Dino	1968-12-26	2012	150
11	A Rella	1989-11-26	2016	20

Visits

visit_	_id member_	_id restaurant_	_id branch_	no	date	score
1	1	4	3		2015-12-30	5
2	8	4	3		2016-01-15	4
3	1	2	1		2018-06-16	5
4	2	3	2		2018-06-14	5
5	1	2	1		2018-12-09	3
6	3	3	2		2018-01-16	6
7	2	4	1		2018-03-26	3
8	4	3	2		2018-08-02	7
9	5	3	1		2018-04-22	8
10	4	5	2		2018-05-10	0
11	5	5	2		2018-05-08	6
12	3	5	1		2018-06-21	1
13	6	5	1		2018-03-06	2
14	4	5	1		2018-02-12	3
15	7	4	2		2018-06-07	1
16	5	4	2		2018-12-04	2
17	3	4	2		2018-01-19	3
18	5	3	3		2018-01-27	5

Branch

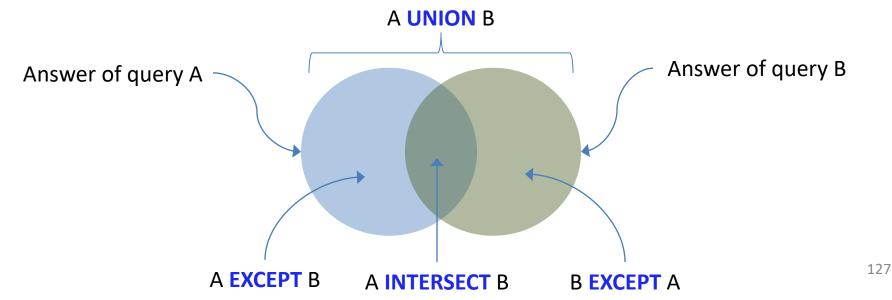
restaurant_	id branch_no	location	seats
1	1	Admiralty	10
1	2	Central	20
2	1	Causeway Bay	5
3	1	Admiralty	25
3	2	Wan Chai	45
3	3	Causeway Bay	35
4	1	Central	170
4	2	Admiralty	100
4	3	North Point	120
5	1	Central	80
5	2	Wan Chai	40

Section 11

Set Operations

Set operations

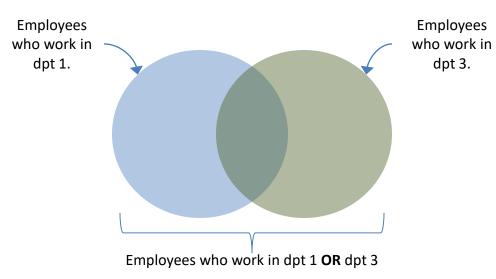
- Set operations can be expressed in SQL using clauses UNION, INTERSECT, EXCEPT.
- Using the set operations can ease the design of SQL by breaking down a complex query to a number of simpler sub-queries.



The UNION clause

Query: Find the names of employees who work in department 1 or department 3.

dpt 3.



Note: The two SQLs are NOT equivalent to each other! **Duplicates are eliminated when** two sets are unified.

SELECT E.name **FROM** Employees E, Works in W WHERE E.employee id= W.employee id AND (W. department id = 1 OR W. department id = 3)

```
SELECT E.name
FROM Employees E, Works in W
WHERE E.employee id = W.employee id AND
      W.department id = 1
```

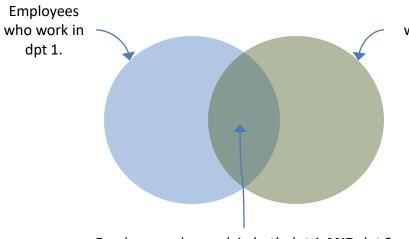
UNION

SELECT E.name **FROM** Employees E, Works in W WHERE E.employee id = W.employee id AND W. department id = 3



The INTERSECT clause

Query: Find the name of employees who work in department 1 and department 3.



Employees who work in dpt 3.

Employees who work in both dpt 1 AND dpt 3

Note: MySQL doesn't support the keyword INTERSECT.

But we can replace **INTERSECT** by joining tables.

SELECT E.name

FROM Employees E, Works_in W
WHERE E.employee_id = W.employee_id AND
W.department id = 1

INTERSECT

FROM Employees E, Works_in W
WHERE E.employee_id = W.employee_id AND
W. department_id = 3

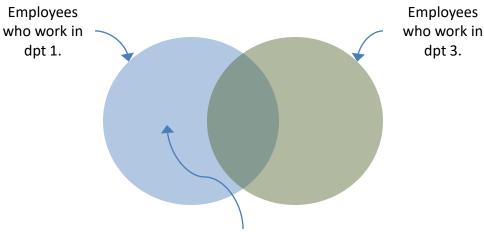


The EXCEPT clause

Query: Find the name of employees who work in department 1 but not department 3.

Employees

dpt 3.



Employees who work in dpt 1 but not in dpt 3

Note: MySQL doesn't support the keyword EXCEPT.

But we can replace **EXCEPT** by using **NOT IN**.

SELECT E.name

FROM Employees E, Works in W WHERE E.employee id = W.employee id AND W.department id = 1

EXCEPT

SELECT E.name **FROM** Employees E, Works in W WHERE E.employee id = W.employee id AND W. department id = 3



Section 12

More on Nested Queries

Nested queries

- Nested queries have other subqueries embedded in them.
- Used for the ease of expressing a natural language request in SQL.
- Subqueries are usually nested under WHERE clauses.
 - May also be enclosed under FROM or HAVING clauses

The N clause

Query: Find the names of the employees in department 1.

```
SELECT E.name
FROM Employees E, Works_in W
WHERE E.employee_id = W.employee_id AND
W.department_id = 1;
```

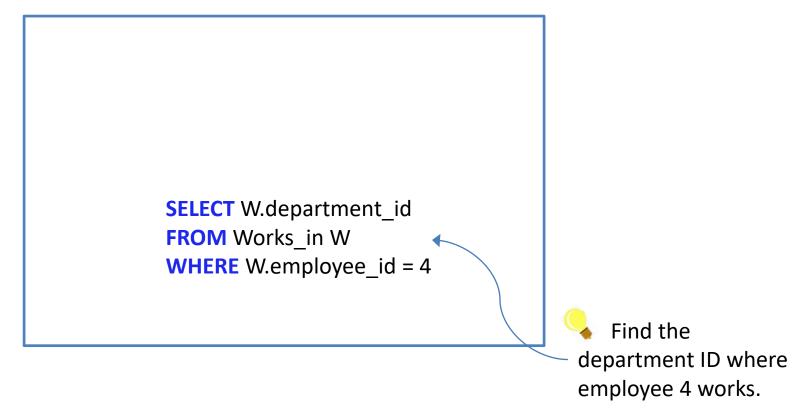
In natural language: Find employee names whose employee_id appears in the set of employee_ids working for department 1.

```
SELECT E.name
FROM Employees E
WHERE E.employee_id IN (
SELECT W.employee_id
FROM Works_in W
WHERE W.department_id = 1);
```

Just like searching the employee_id in the result of the nested query.

The **SOME** clause

Query: Find department names that have greater budget than some department where employee 4 works.



The **SOME** clause

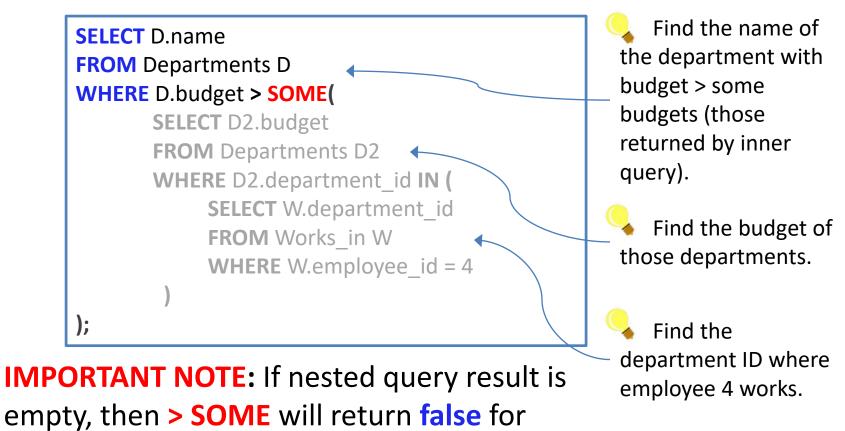
Query: Find department names that have greater budget than some department where employee 4 works.

```
SELECT D2.budget
FROM Departments D2
WHERE D2.department id IN (
     SELECT W.department id
                                              Find the budget of
     FROM Works in W
                                           those departments.
     WHERE W.employee_id = 4
                                              Find the
                                           department ID where
                                           employee 4 works.
```

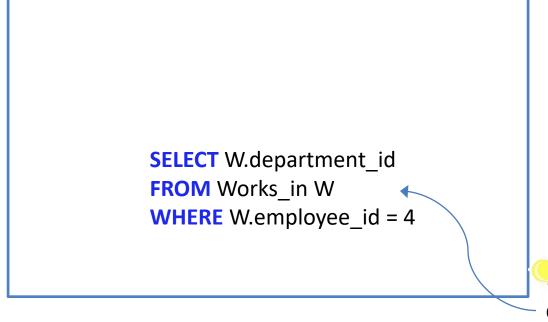
The **SOME** clause

every D.budget!

Query: Find department names that have greater budget than some department where employee 4 works.



Query: Find department names that have greater budget than all departments where employee 4 works.



Find the department ID where employee 4 works.

Query: Find department names that have greater budget than all departments where employee 4 works.

```
SELECT D2.budget
FROM Departments D2
WHERE D2.department_id IN (
SELECT W.department_id
FROM Works_in W
WHERE W.employee_id = 4
)

Find the budget of those departments.
```

Query: Find department names that have greater budget than all departments where employee 4 works.

```
SELECT D.name
FROM Departments D
WHERE D.budget > ALL (
SELECT D2.budget
FROM Departments D2
WHERE D2.department_id IN (
SELECT W.department_id
FROM Works_in W
WHERE W.employee_id = 4
)
);
```

Find the name of the department with budget > ALL budgets (those returned by inner query).

IMPORTANT NOTE: If nested query result is empty, then
> ALL will return true for every D.budget!

Query: Find department names that have the greatest budget than all departments.

Question: What would the result be if >ALL is used?



Can you rewrite the above query using **Aggregate function MAX** in a **nested query**?

The EXISTS clause

- The inner subquery could depend on the row currently examined in the outer query.
 - Query: Find the names of employees who work in department with department_id=1.

```
SELECT E.name
FROM Employees E
WHERE EXISTS (
SELECT *
FROM Works_in W
WHERE W.department_id = 1 AND
E.employee id = W.employee_id);
For each employee record r.
If the inner query can return some records, we will return the record r.
```

EXISTS is a boolean set-comparison operator that returns false if the input set is empty and true otherwise.

Section 13

Null values

NULL value

- Handling null values is a non-trivial topic in database research.
- null: unknown value or value does not exist.
- Use predicate IS NULL to check for null values.
 - Query: Find all employee names for which the salary is unknown or undetermined.

Employees

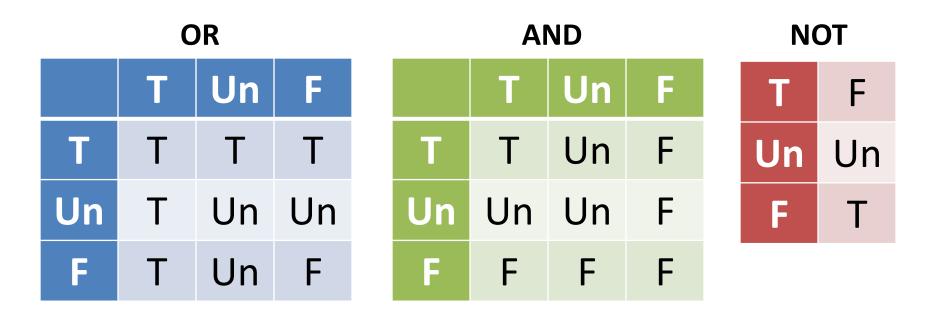
employee_id	name	salary
1	Jones	
2	Smith	28000
3	Parker	
4	Smith	24000

SELECT name
FROM Employees
WHERE salary IS NULL

NULL value

- The result of any arithmetic expression involving null is null.
 - 5 + null returns null.
- Any comparison with null returns UNKNOWN.
 - Both 5 < null, null = null return UNKNOWN.</p>
- Use P IS UNKNOWN to check if a predicate P is unknown or not.
- For the result of WHERE or HAVING clause, predicate is false if it evaluates to UNKNOWN.

Three valued logic



NULL value and aggregates

SELECT SUM (budget) **FROM** Departments

- The statement above ignores null amounts.
- All aggregate operations except COUNT(*) ignore tuples with null values on the aggregated attributes.
 - COUNT counts not null values only.

SUM(budget) returns 100000.

COUNT(*) returns 3.

COUNT(budget) returns 1.

department_id	name	budget
1	Toys	
2	Tools	
3	Food	100000

Section 14

Views

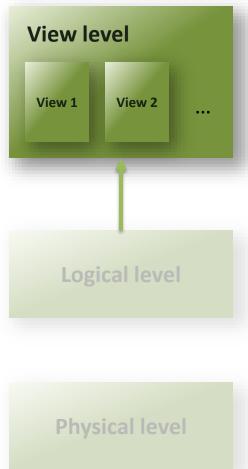
The CREATE VIEW clause

Views provide a mechanism to hide certain data from the view of certain users.

Syntax:

CREATE VIEW view_name AS <expression>

```
CREATE VIEW Employee_hide_salary AS (
SELECT employee_id, name
FROM Employees);
```



Section 15

Authorization

Authorization

The DBA can grant access/update authorization to users.

Syntax:

GRANT <priviledge list>

ON

TO <user/role list>

Johnson, Brown are usernames.

GRANT SELECT ON Departments **TO** Johnson, Brown

GRANT UPDATE(budget) **ON** Departments **TO** Johnson

GRANT UPDATE(budget) **ON** Departments **TO** manager

manager is not a username, it is a **role**.

Authorization

Rights can be revoked.

REVOKE SELECT ON Departments **FROM** Johnson, Brown

Create a role.

CREATE ROLE manager;

Grant a role to a user.

GRANT manager **TO** Brown;

Section 16

Assertion

Assertions

- An assertion ensures a certain condition will always exist in the database.
 - Assume that we want to enforce that the number of departments cannot exceed the number of employees at any valid instance of our database.

Assertions are checked every time the involved tables are updated and they could be very expensive.

A common table expression (CTE) is a named temporary result set that exists within the scope of a single statement and that can be referred to later within that statement, possibly multiple times.

```
WITH cte1 AS (SELECT a, b FROM table1),
cte2 AS (SELECT c, d FROM table2)
SELECT b, d FROM cte1
INNER JOIN cte2
ON cte1.a = cte2.c;
```

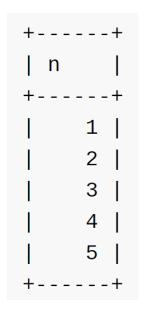
```
WITH cte (col1, col2) AS
(SELECT 1, 2
UNION ALL
SELECT 3, 4)
SELECT col1, col2 FROM cte;
```

```
WITH cte (col1, col2) AS
(SELECT 1, 2
UNION ALL
SELECT 3, 4)
SELECT col1, col2 FROM cte;
```

```
col1 col21 23 4
```

A recursive common table expression is one having a subquery that refers to its own name.

```
WITH RECURSIVE cte (n) AS
( SELECT 1
UNION ALL
SELECT n + 1 FROM cte
WHERE n < 5 )
SELECT * FROM cte;
```



The first <u>SELECT</u> produces the initial row or rows for the CTE and does not refer to the CTE name. The second <u>SELECT</u> produces additional rows and recurses by referring to the CTE name in its <u>FROM</u> clause. Recursion ends when this part produces no new rows. Thus, a recursive CTE consists of a nonrecursive <u>SELECT</u> part followed by a recursive <u>SELECT</u> part.

The recursive CTE subquery shown earlier has this nonrecursive part that retrieves a single row to produce the initial row set:

```
1 SELECT 1
```

The CTE subquery also has this recursive part:

```
1 SELECT n + 1 FROM cte WHERE n < 5
```

At each iteration, that $\underline{\mathtt{SELECT}}$ produces a row with a new value one greater than the value of n from the previous row set. The first iteration operates on the initial row set (1) and produces 1+1=2; the second iteration operates on the first iteration's row set (2) and produces 2+1=3; and so forth. This continues until recursion ends, which occurs when n is no longer less than 5.

A Fibonacci series begins with the two numbers 0 and 1 (or 1 and 1) and each number after that is the sum of the previous two numbers.

```
WITH RECURSIVE fibonacci (n, fib_n, next_fib_n) AS
(SELECT 1, 0, 1
UNION ALL
SELECT n + 1, next_fib_n, fib_n + next_fib_n FROM fibonacci
WHERE n < 10)
SELECT * FROM fibonacci;
```

+	-+	-++
n	fib_n	next_fib_n
+	-+	-++
1	0	1
2	1	1
3	1	2
4	2	3
5	3	5
6	5	8
7	8	13
8	13	21
9	21	34
10	34	55
++		

Lecture 4

END

COMP3278A

Introduction to Database Management Systems

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